

## CONVERGENCE THEORY

- **Background: Best uniform approximation;**
- **Chebyshev polynomials;**
- **Analysis of the CG algorithm;**
- **Analysis in the non-Hermitian case (short)**

### Background: Best uniform approximation

We seek a function  $\phi$  (e.g. polynomial) which deviates as little as possible from  $f$  in the sense of the  $\|\cdot\|_\infty$ -norm, i.e., we seek the

$$\min_{\phi} \max_{t \in [a,b]} |f(t) - \phi(t)| = \min_{\phi} \|f - \phi\|_\infty$$

where  $\phi$  is in a finite dimensional space (e.g., space of polynomials of degree  $\leq n$ )

- Solution is the “best uniform approximation to  $f$ ”
- Important case:  $\phi$  is a polynomial of degree  $\leq n$
- In this case  $\phi$  belongs to  $\mathbb{P}_n$

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Text: 6.11 – cheby

### The Min-Max Problem:

$$\rho_n(f) = \min_{p \in \mathbb{P}_n} \max_{x \in [a,b]} |f(x) - p(x)|$$

- If  $f$  is continuous, best approximation to  $f$  on  $[a, b]$  by polynomials of degree  $\leq n$  exists and is unique
- ... and  $\lim_{n \rightarrow \infty} \rho_n(f) = 0$  (Weierstrass theorem).

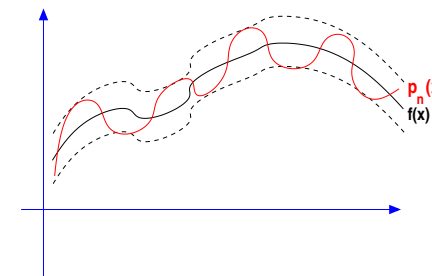
**Question:** How to find the best polynomial?

**Answer:** Chebyshev’s equi-oscillation theorem.

Chebyshev equi-oscillation theorem:  $p_n$  is the best uniform approximation to  $f$  in  $[a, b]$  if and only if there are  $n + 2$  points  $t_0 < t_1 < \dots < t_{n+1}$  in  $[a, b]$  such that

$$f(t_j) - p_n(t_j) = c(-1)^j \|f - p_n\|_\infty \quad \text{with } c = \pm 1$$

[ $p_n$  ‘equi-oscillates’  $n + 2$  times around  $f$  ]



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### Application: Chebyshev polynomials

**Question:** Among all monic polynomials of degree  $n + 1$  which one minimizes the infinity norm? Problem:

$$\text{Minimize } \|t^{n+1} - a_n t^n - a_{n-1} t^{n-1} - \dots - a_0\|_\infty$$

**Reformulation:** Find the best uniform approximation to  $t^{n+1}$  by polynomials  $p$  of degree  $\leq n$ .

➤  $t^{n+1} - p(t)$  should be a polynomial of degree  $n + 1$  which equi-oscillates  $n + 2$  times.

➤ Define Chebyshev polynomials:

$$C_k(t) = \cos(k \cos^{-1} t) \text{ for } k = 0, 1, \dots, \text{ and } t \in [-1, 1]$$

➤ Observation:  $C_k$  is a polynomial of degree  $k$ , because:

➤ the  $C_k$ 's satisfy the three-term recurrence :

$$C_{k+1}(t) = 2tC_k(t) - C_{k-1}(t)$$

with  $C_0(t) = 1, C_1(t) = t$ .

 Show the above recurrence relation

 Compute  $C_2, C_3, \dots, C_8$

 Show that for  $|t| > 1$  we have

$$C_k(t) = \text{ch}(k \text{ ch}^{-1}(t))$$

- $C_k$  Equi-Oscillates  $k + 1$  times around zero.
- Normalize  $C_{n+1}$  so that leading coefficient is 1

The minimum of  $\|t^{n+1} - p(t)\|_\infty$  over  $p \in \mathbb{P}_n$  is achieved when  $t^{n+1} - p(t) = \frac{1}{2^n} C_{n+1}(t)$ .

➤ Another important result:

Let  $[\alpha, \beta]$  be a non-empty interval in  $\mathbb{R}$  and let  $\gamma$  be any real scalar outside the interval  $[\alpha, \beta]$ . Then the minimum

$$\min_{p \in \mathbb{P}_k, p(\gamma)=1} \max_{t \in [\alpha, \beta]} |p(t)|$$

is reached by the polynomial:  $\hat{C}_k(t) \equiv \frac{C_k\left(1 + 2\frac{\alpha-t}{\beta-\alpha}\right)}{C_k\left(1 + 2\frac{\alpha-\gamma}{\beta-\alpha}\right)}$ .

### Convergence Theory for CG

➤ Approximation of the form  $x = x_0 + p_{m-1}(A)r_0$ . with  $x_0 =$  initial guess,  $r_0 = b - Ax_0$ ;

➤ Recall property:  $x_m$  minimizes  $\|x - x_*\|_A$  over  $x_0 + K_m$

➤ **Consequence:** Standard result

Let  $x_m = m$ -th CG iterate,  $x_*$  = exact solution and

$$\eta = \frac{\lambda_{\min}}{\lambda_{\max} - \lambda_{\min}}$$

$$\text{Then: } \|x_* - x_m\|_A \leq \frac{\|x_* - x_0\|_A}{C_m(1 + 2\eta)}$$

where  $C_m =$  Chebyshev polynomial of degree  $m$ .

- Alternative expression. From  $C_k = ch(kch^{-1}(t))$ :

$$C_m(t) = \frac{1}{2} \left[ \left( t + \sqrt{t^2 - 1} \right)^m + \left( t + \sqrt{t^2 - 1} \right)^{-m} \right]$$

$$\geq \frac{1}{2} \left( t + \sqrt{t^2 - 1} \right)^m. \quad \text{Then:}$$

$$C_m(1 + 2\eta) \geq \frac{1}{2} \left( 1 + 2\eta + \sqrt{(1 + 2\eta)^2 - 1} \right)^m$$

$$\geq \frac{1}{2} \left( 1 + 2\eta + 2\sqrt{\eta(\eta + 1)} \right)^m.$$

- Next notice that:

$$1 + 2\eta + 2\sqrt{\eta(\eta + 1)} = \left( \sqrt{\eta} + \sqrt{\eta + 1} \right)^2$$

$$= \frac{(\sqrt{\lambda_{\min}} + \sqrt{\lambda_{\max}})^2}{\lambda_{\max} - \lambda_{\min}}$$

$$= \frac{\sqrt{\lambda_{\max}} + \sqrt{\lambda_{\min}}}{\sqrt{\lambda_{\max}} - \sqrt{\lambda_{\min}}}$$

$$= \frac{\sqrt{\kappa} + 1}{\sqrt{\kappa} - 1}$$

where  $\kappa = \kappa_2(A) = \lambda_{\max}/\lambda_{\min}$ .

- Substituting this in previous result yields

$$\|x_* - x_m\|_A \leq 2 \left[ \frac{\sqrt{\kappa} - 1}{\sqrt{\kappa} + 1} \right]^m \|x_* - x_0\|_A.$$

- Compare with steepest descent!

### Theory for Nonhermitian case

- Much more difficult!
- No convincing results on ‘global convergence’ for most algorithms: FOM, GMRES(k), BiCG (to be seen) etc..
- Can get a general a-priori – a-posteriori error bound

### Convergence results for nonsymmetric case

- Methods based on minimum residual better understood.
- If  $(A + A^T)$  is positive definite ( $(Ax, x) > 0 \forall x \neq 0$ ), all minimum residual-type methods (ORTHOMIN, ORTHODIR, GCR, GMRES,...), + their restarted and truncated versions, converge.
- Convergence results based on comparison with one-dim. MR [Eisenstat, Elman, Schultz 1982] → not sharp.

MR-type methods: if  $A = X\Lambda X^{-1}$ ,  $\Lambda$  diagonal, then

$$\|b - Ax_m\|_2 \leq \text{Cond}_2(X) \min_{p \in \mathcal{P}_{m-1}, p(0)=1} \max_{\lambda \in \Lambda(A)} |p(\lambda)|$$

(  $\mathcal{P}_{m-1} \equiv$  set of polynomials of degree  $\leq m - 1$ ,  $\Lambda(A) \equiv$  spectrum of  $A$  )