Csci 4061 - Meeting 1

• Administrative

- * Introductions
- * Syllabus
- * Waiting List Policy
- Goals:
 - * Understand concurrency
 - * Overview of the course
- Topics:
 - * 1.1 Multiprogramming and Multitasking
 - * 1.2 Concurrency at the Application Level
 - * 1.3 Unix Standards
 - * 1.4 Programming in UNIX
 - * 1.5 Making functions safe
- *Readings: Chapter 1 (Robbins, pp.76-137)*
- Readings: Appendix A.1-3 (Robbins, pp.577-589)
- Recommended Exercises:

Administrative

- Introductions
 - * Instructor
 - * TAs
 - * Peers
- Syllabus
 - * Schedules: lectures, homeworks, exams, recitations
 - * Textbooks, Reference material
 - * Topics
- Policies
 - * Late homeworks
 - * Cheating
 - * Waiting List Attendance*, estimates

Course Goals

- Concepts: Understand concurrency
 - * Why concurrency?
 - * Sources of Concurrency
 - I/O, signals, processes, threads, client-server
 - * Effects of concurrency
- Focus
 - * Server software concurrently shared by many
 - * User level commands, shell
 - * Power Users system calls, C programs
- Out of Scope
 - * Operating System Theory e.g. CPU scheduling
 - * Vendor specific features e.g. Win32

1. What is Concurrency?

• Concurrency:

- * Sharing of resource in the same time-frame
- * Ex. two program executing concurrently
- * Q? Which resources are they sharing?

• Trends leading to Concurrency

- * Computer speed >> Human typing speed
- * CPU speed >> I/O (e.g. disk drives)
- See Table 1.1 (pp. 5)
- * Multiprocessors
- * Distributed Systems
- * Graphical User Interfaces
- Animation of multiple objects
- What is hard about Concurrency?
 - * Non-deterministic behaviour
 - * Bugs do not show up on a regular basis

1.1 Multiprogramming and Multitasking

• Multiprogramming

- * Process: instance of a program in execution
- * More than one processes can be ready to execute
- * OS chooses one to execute
- * Context switch to another process when
- this process needs I/O
- Q? What if a program has an infinite loop?
- Timesharing
 - * Context switch to another process when
 - when quantum is over
 - * Pros: Reduce waiting time for small jobs
 - * Cons: overhead of context switch
- Multitasking Similar to multiprogramming
 - * finer granualarity (e.g. threads within a process)
 - * Sharing even user resource, e.g. global variables
- Why do I care about these? I am not writing an O.S.!
 - * Web servers: search engines, databases, e-commerce

1.2 Concurrency at the Application Level

• Concurrency Levels

- * Hardware
- CPU controlling peripherals, multi-processors
- * Software OS
- signal handling
- overlap of I/O and processing
- communication
- resource sharing among processes and threads

• Outline

- * 1.2.1 Interrupts (Recall assembly language course)
- * 1.2.2 Signals
- * 1.2.3 Input and Output
- * 1.2.4 Threads and Resource Sharing
- * 1.2.5 Network as the Computer
- Q? Map these to chapters in the book.

1.2.1 Interrupts

- * Peripheral generates an electrical signal
- * Sets a flag in CPU
- * CPU checks flag in each instruction cycle
- * Interrupt service routine called
- Example: Timesharing implemented with
 - * alarm interrupts
- Concurrency: CPU and peripheral device
 - * Shared resource bus
- Event types by time of occurrence
 - * Asynchrounous time not determined by receiver
 - * Synchronous time determined by receiver

1.2.2 Signals

• Motivation

- * Q? How do you stop a program in an infinite loop?
- * Other usage: timers, job control, aynch. I/O, ...
- Signal = software notification of an event
 - * Ex. hardware events, e.g. ctrl-c, I/O complete
 - * Q? Provide examples of synchronous signals.
- Life cycle of a Signal
 - * Event of interest occurs
 - * Signal is generated
 - * OS sets a flag for the relevant process
 - * Signal is caught by the process
 - * Process invokes a handler subroutine
 - * Analogy "You have mail" flag
- Concurrency: main program, signal handler subroutine
 - * Implication: restriction on signal handler
 - * Sharing a global variable => special protection

1.2.3 Input and Output

• Motivation

- * Coordinate resources with varying speed
- * But isn't that the job of O.S.?
- * Why should an application developer learn this?
- * You may develop performance critical applications
- Ex. real-time Pacemaker
- Ex. Web servers, transaction processors ebay, amazon, ...
- Ex. asynchronous I/O
 - * A process itseld can do other things
 - * while waiting for an I/O, i.e. synchronous read()
 - * instead of getting swapped out by OS
- Ex. monitoring multiple input source on network
 - * Standard blocking I/O is not suitable!
- Concurrency
 - * Subprogram handling file/network I/O
 - * Subprograms computing during wait for I/O

1.2.4 Threads and Resource Sharing

- Motivation What is the unit of concurrency?
 - * Traditional unit = process
 - * Emerging finer unit = thread
- Processes Generated via fork() call
 - * Coordinate termination via wait()
 - * Communicate via pipes (common ancestors),
 - or signals, messages, shared memory, etc.
 - * Pros: stronger security boundaries
 - * Cons: high overhead
- Threads provide concurrency within a process
 - * threads of execution = program counter value streams
 - * Finer level of concurrency
 - * Low overhead in creating and context switching
 - * standards are emerging now!
- Concurrency
 - * Multiple processes or Multiple threads within a process

1.2.5 Network as the Computer

- Motivation internet!, intranet, networks, ...
 - * Multiple services: ftp, email, ...
 - * Million of clients accessing Web services
- *Client-Server* = A model of distributed computing
 - * Client = caller of a service
 - * Server = provider of a service
 - * Analogy with procedure call, caller, callee
- Details
 - * Clients and Servers may be on different machines
 - * Communication via messages or remote procedure calls
 - * Signals, Pipes, shared memory are not common
- Concurrency
 - * Server and client are concurrent
 - * Multiple Servers and multiple clients

1.3 Unix Standards

• Why Standards?

- * Multiple flavours of Unix: HPUX, Solaris, Linux, ...
- Two distinct lineage BSD and System V
- * Non-Unix OS: NT, Windows 3.1/95/98/..., MacOS, ...
- * System calls are often OS specific!
- * Overhead of porting across OS.
- Which Standards?
 - * ANSI C
 - * POSIX IEEE Portable Operatig System Interface
 - Table 1.3 provide POSIX standards
 - * if not covered by POSIX
 - Spec 1170
 - System V Release 4
- How do I check POSIX support in my OS?
 - * unistd.h header file
 - * Table 1.4 shows the compile time options

1.4 (Concurrency) Programming in UNIX

- Concurrency programming
 - * Language constructs, e.g. Java
 - * OS Libraries, e.g. Unix system call
- System call a procedure provided by OS
 - * An entry into the kernel (heart) of OS
 - * To get access to system resources
- Standard C Library
 - * e.g. string handling, memory management
 - * Some subroutine contain system calls
 - * Hard to tell the difference from system call!
- Resources Unix man pages (Appendix A.1)
 - * header files needed by system call
 - * prototype of system call- name, parameters
- Apendix A.1-3commands: man, cc, make

1.4 Programming in UNIX

• Conventions- error situation

- * system call returns -1 or NULL
- Sets global variable "errno" to error code
- * Application programmer should check for these
- perror() Example 1.2 (pp. 15)
- strerror() Example 1.5 (pp.16)
- Newer Style use exception handling (C++, Java)
- new system calls return error code as result
- avoid global variables, e.g. "errorno"
- Other conventions (See bullets on pp. 16-17)
 - * Q? Identify 3 bullet related to concurrency?
 - * Q? Which bullet relates to memory leaks?
 - * Q? List problems with global variable "errno".

1.4 Programming in UNIX

- Extended example argument arrays!
 - * Review pointers, argv[], argc, parameter passing
- *Ex. Review Program 1.1 and 1.2 to answer the following:*
 - * What are argv[] and argc used for?
 - * What is the parameter passing mode in C?
 - * What are the data types of arguments to makeargv()?
 - * What does makeargv() return?
 - * List a few possible error situations for makeargv().
 - How does makeargv() respond to those errors?
 - * Is it possible to rewrite makeargy() with following header?
 - Headers from Example 1.8, Example 1.10
 int makeargv(char *s, char *delimiters, char **argvp)
 - * What is maximum number of arguments allowed?
 - * Is there any memory leak? Justify your answer.
 - Consider memory allocated to 't' and '*argvp'
 - * What the following loop do?

for (i=1; i < numtokens + 1; i++)

*((*argvp) + i) = strtok(NULL, delimiters);

* Why is the above loop not followed by free(t)?

1.5 Making functions safe (for reentry)

• Non-Reentrant functions

- * Self modifying code
- * functions using static/global variables
- * Problems with multiple simultaneous invocations

• Reentrant functions

- * Allow multiple simultaneous invocations
- * Needed for signal handler, server with many clients, ...
- * Two aspects -
- Thread safe: can be called concurrently by 2 threads
- Async. Signal safe: can be called inside a signal handler
- without restriction

1.5 Making functions safe (for reentry)

- Q? Which POSIX system calls thread safe?
 - * Not those using global variable "errno", e.g. read()
 - * reentrant functions provided for non-reentrant ones
 - * Ex. strtok_r() for strtok()
 - * Trend towards thread safe system calls!
- Q? Which POSIX system calls async signal safe?
 - * See Table 5.3, pp. 191
 - * Double check with man page on your system!
- Q? Is makeargv() (Program 1.2) a reentrant function?
 - * Is it signal safe? Is it thread safe? Why?
 - * How can you make it thread safe?